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Privacy-friendly Platform for Healthcare Dat. ir Cloud Based on Blockchain Environment

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Abstract

Data in cloud has always been a point of traction for the cyber attackers. Nowadays healthcare data in cloud has become their new interest. Attacks on these healthcare data can result in annihilating consequences for the healthcare organizations. Decentralization of these cloud data can minimize the effect of attacks. Storing and running computation on sensitive private healthcare data in cloud are possible by decentralization which is enabled by peer to peer (P2P) network. By leveraging the decentral red or distributed property, blockchain technology ensures the accountability and regrity. Different solutions have been proposed to control the effect of attacks using decentralized approach but these solutions

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somehow failed to ensure overall privacy of patient centric system. Ir this paper, we present a patient centric healthcare data management system $l_{ing} l_{i}^{i}$ ckchain technology as storage which helps to attain privacy. Cryptographic tull ctions are used to encrypt patients data and to ensure pseudonymity. We many ze the data processing procedures and also the cost effectiveness of the snumt contracts used in our system.

Keywords: Blockchain, Decentralization, Healthcare data in cloud, Pseudonymity, Privacy, Security, Smart contract

1. Introduction

A lot of work is going on healthcare and intermation technology in an amalgamated manner and these works are bring the intermediate of changes in healthcare discipline. These changes are affecting patients' the attent process hence requiring careful data processing. For treatment, health are is completely dependent on data which arises some concerns over data the curity and privacy. Authorization or private access to the personal data of in the dual patient refers to the term Privacy, which means only authenticated there will be able to access the private data. Keeping these personal data safe from the eavesdroppers or intruders refers to the term Security, which means system will be able to protect users' private data from outsiders. Authenticated parties of healthcare data preservation process will get



Figure 1: Entities of EHR system and it's Data flow

the access to store data into cloud and retrieve from it. Interaction between the system and the patient requires a secured channel. Different authoritical on protocol [20, 19, 24] have been proposed to preserve the privacy ar a security. Lack of security may result in devastating consequences like data los' and data theft. A lot of intruders are searching for an insecured channel and trying access valuable healthcare data in the cloud network. In most of the case , data 'oss in healthcare causes detrimental consequences to the patients and health care or ganizations. Due to recent attacks on healthcare data in cloud systems, inferent countries like USA [8] and UK [12] have experienced critical data los. resc nal data of patients' were kept without encryption in the cloud which all wed ".e attackers to steal the sensitive private data. Let's assume a scenario where the tients keep their data in any Electronic Health Record (EHR) system [35, 1, 14, 13, 38, 5] for preservation and also for further access. Figure 1 depicts a gen. ralized formation of EHR systems. In the figure patients and healthcare of minizations take part in the process as both data sender and data receiver. EHR system : the manager of the whole process that maintains the data flow of the system. Top most entity is the cloud where data is kept. Patients share their perscal data with the doctors and healthcare organizations with the help of these EHR systems. Suppose, a patient keeps her data in the cloud system [7] which uses we ain as a data storage platform. System will store the data on blockchain when $d \circ patient$ shares her data with the system. Accountability of data is system centric in case of the instance [7], whereby the system will provide data stor ge serv ce even when data is shared with the doctors or healthcare organizations Con or lently, the system is responsible for data loss.



Figure 2^1 depicts the design of our platform in which aforen. Intioned problems have been addressed by storing the encrypted healthcare using in the cloud system. As a result, if our system somehow loses the contrainver elockchain, patients will be accountable for their data as they will contrain the encryption keys solely. Data sharing in our system is also being controlled by the patients. Vulnerabilities related to data preservation have been addressed in our system by using cryptographic functions along with blockchain technology. How ever, our system will store the encrypted personal data ensuring overall privacy of the data such that even if system gets attacked by the attacker the stole radata vill make no sense to them. To get the plaintext of those encrypted personal data vill make no sense to the keys. There is no identifier for these datasets, only encryption keys will be used to identify such encrypted and pseudonymou.² data.

1.1. Our Contribution

Our platform ensures that the private e^{-1} heare data in cloud is controlled by only patient herself. The main idea of this work is to keep the sensitive healthcare data on the blockchain to attain account billty, integrity and security. Patients will have the overall control over the block in which their data will be stored. Present healthcare systems lack in pseudonymic, as those only store the data in cloud, but our platform ensures the pseudonymic as those only store the data in cloud, but our platform ensures the pseudonymity of patients. We achieve pseudonymity by using cryptographic function. Med. Behain will regain the interest of patients on EHR systems and will retain acourtability, integrity, pseudonymity, security and privacy which are being left with the increasing computational power of emerging technologies in EHR systems.³ Analysis of these attributes is discussed in section 3. Our contributions e^{-e} as follows:

- 1. Security and rivery guarantee: The proposed platform guarantees accountability ps udonymity, authenticity and integrity along with data privacy.
- 2. Analys... 'Ligo ous analysis on security, privacy, accountability, pseudonymity and 'Legrity nows how our platform achieves the above mentioned properti s.

¹P₁ vate Accessible Unit (PAU) is the intermediary unit between blockchain and data sender or receive.

²Pseudonymity refers to the fact of using disguised identity.

An aysis of security terminologies are given in section 5.

3. Evaluation: We have implemented smart contract and show. different analogies of costs (e.g., transaction cost, execution cost). Then we have evaluated a java implementation of input and output generation algorithm using Elliptic Curve Cryptography (ECC) for our system. Experimental results will help to compare several aspects of EHR system and will help to decide whether accept our platform or not.

Organization of the paper: The remainder of the paper is organized as follows: Section 2 describes the related work. In section 3^{v} ed scuss the preliminaries. In section 4, we describe our platform. In Section 5, we evaluate the platform and analyze it formally. We give some concluding remarks in Section 6.

2. Related Work

Some national level frameworks based on civil for electronic medical system have been proposed in [14, 13, 25]. Para e 1 [25] proposed a model which is cloud-based and deals with patients' private data. This model ensures cost effectiveness, and this system was designed for rural areas where cost plays an immense role. Medical profession loand policy makers could serve the patients remotely through a cloud-based mood' which stores all the imperative data in a single cloud. Patients were encouraged to share their data in the cloud so that they could get the medical service i om the professionals remotely. Disease diagnosis and control could be n_{1} de l_{y} this remote treatment. Data collection and data delivery are the key r sint in symptom analysis. Rolim et al. [27] proposed a framework where the system provesses data in the steps of data collection and data delivery. In this mode sensor, play the role of collector which collects the data and sends directly to the system to store and work with this data further. These data would be accossed by the medical professionals and sensors were proposed to be attached with the medical equipment in this system. Yin et al. [37] introduced cloud based particentric system. This model includes three layers: data collection lay *i*r, *c* at a management layer and data service layer. [21] described a blockchain base. ac ess control manager for heath data to enhance the interoperability of nis sy tem. Off blockchain mechanism with the involvement of public blockchai. was p oposed as an access control manager of healthcare data.

Controlla pility and Traceability are two key topics of privacy preserving systems. All all [35] proposed a model which is based on blockchain to help patients where the personal data easily and securely with privacy

preservation. This application based model also deals with Sec. ~ / Aulti-party Computing (MPC) and Indicator-Centric Schema (ICS). Simic et 1. [2] showed a case study where the study concludes with the illustration of an inficant benefits of IoT and blockchain in a combined manner. In their work I/T devices were proposed to be used as collectors of private health data of the patients', and real time data of patient could be saved in blockchain. Scal bility of the blockchain in case of Big data has also been tested in their study. Ek'law e al. [7] proposed a prototype named 'MedRec' which uses blockchair as a backbone and tried to find the security solutions for EHR systems. They rie to give their prototypeintegrity, authenticity, auditability and data sharing "brouz" olockchain. Elements of their system are: Registrar Contract (RC), Patient-Provider Relationship Contract (PPR), Summary Contract (SC), where RC n. ps t' e identification strings of the participants to their Ethereum addresses, PK. issues contracts between two nodes in the system when one node stores a manages medical records for the other, SC locates the participants medical record history. Jun et al. [3] proposed a web-based architecture where they shower a secured accessing multiple patient repository system. They concent ted h ainly on lifetime repository of health data, which consists of client applicatio. (CA), central access-control (CAC), local access-control (LAC) and Hospital information system. Linn et al. [21] described a blockchain based access conul manager for health data to enhance the interoperability of this system

The backbone of our work ... b' sckchain. Blockchain technology is popular for its application in Biter in c yptocurrency [26], which is a public ledger to hold and maintain the transport all ata and integrity [31]. One of the reasons for using blockchain technology in cryptocurrency is its decentralized digital ledger property, which was presented by Nakamoto [22] in his Bitcoin cryptocurrency framework. Blockchain's data structure has been modeled by blocks which is linearly sequence.' Fach block contains the cryptographic hashes corresponding to the previour and current block to ensure continuity and immutability of the chain. Chaining riech inism ensures integrity of this secured data structure.

2.1. Bloc chain.

Figure ? exhibits the structure of blocks in the blockchain network. In the figure eren block is connected to its previous block by the hash of previous block. Blocks store he time-stamp of being mined in the network. Mining takes place in the network by solving mathematically complex problems. Miners compete each



Figure 3: Structure of Blocks in bloc' chain

other to mine the block so that they could earn some propocurrency. In our platform miners will get Ether from Ethereum Network for Lining, and our platforms Ethereum account will be charged against it. Simple Cher transfer functionality will be used to transfer the Ether from our account Each block contains corresponding block number and data that has been given to store in the blockchain which has been denoted as δ_n .

Blockchain-secured transaction-based tec. n logy [1] gives the users a better security. Bitcoin as well as blockchain and been failed since these were introduced [6]. The network is shared and information is stored throughout the whole network, thus increasing the reliab." yo, this technology. All the information is treated in a redundant way in blockchain [28]. Blockchain is distributed but it remains all the same for it' not's ensuring the integrity [4, 34]. Centralized database can be corrupted an, needs third party to maintain it. To change the history of the blockchain any 'ndivid. ' has to control at least 51% of the chain and it will cost a lot to challen ; the impattability of blockchain. This immutable architecture [2, 30, 32] is a blessner in archival science too. Identities in the blockchain are covered by pseud n. ms by which privacy for the participants is ensured with a very high degree [15]. Cyptographic authentication of the time blocks with time-stamp allow the entire network to hold the logs for any interaction in the blockchain. Block, in ensures the verifiability of the users. Other than above discussed characteristics some author explicitly mention the key points like trust enabling notio, [1, 2⁷, 11, 33], Consensus, Transparency, Smart contract etc.

Blockchair gives a distribution oriented service to be used as a storage. All the records that may be stored in the blockchain have to use smart contracts[16, 9]. Smart contracts determine the record of data and conditions in the blockchain. These contracts, as a form of code, give a huge power to the programmers to read and write ov r the blockchain [9]. As storage, blockchain provides accuracy and reliability to it's users and protects the data from fraud and being tampered or

corrupted [18]. Blockchain as storage maintains proper decentral. vi on and true redundancy, total privacy and cost reduction [10]. Decentralized veb vill be the future of this era[36].

3. Preliminaries

In this Section, we explain each properties (e.g., security, r rivacy and management) that our protocol achieves. Finally, we intraface the building blocks of our protocol.

3.1. Properties

3.1.1. Security and Privacy

We briefly describe each of the security and \mathbf{P}^{-1} ivacy properties in the context of our system below.

- 1. Pseudonymity: No entity will be one identify any party of our system because users are being identified by a dynamic key. As a result users are keeping their selves pseudonymou.⁴ Data will not be identified by just seeing it.
- 2. Privacy: Only registered parties will be able to interact with the system. Even a registered party win pot be able to access the private raw data of other parties.
- 3. Integrity: Authenti ater parties will be able to store private data.
- 4. Accountability: Each back will be identified by corresponding block-id. Only authenticated parties will get them and will interact with them.
- 5. Security: P'.rtie' will keep their encrypted data in the system which ensures secured envny ment for them.
- 3.1.2. Mana₈, m. nt
 - User, need to register once and by providing the ID and PWD ⁵ they can eas 'y get it to the platform.

⁴Pe sudonyn ty and anonymity are two different things. Anonymity refers to the fact of being unknow. in $\circ r$ system users are identified with dynamic keys, hence users are pseudonymous. Sim and PWD are described in Table 1.

- PAU will act as a Trusted Third Party (**TTP**) of our system, . it will be the medium between user and blockchain.
- In the case of Block id sharing, users need to be very car (... because untrusted access will make the platform vulnerable for u.) particular user's data.

3.2. Cryptographic tools

Here, we describe Elliptic Curve Cryptography (ECC) [17] which has been used as the cryptographic tool to provide proper cryptographic functionality to the users. Formal definition of ECC will be given here.

Definition 1 (Elliptic Curve Cryptography) Enliptic Curve Cryptographic scheme use the trapdoor function which means if w compute B from A through trapdoor function then it is mathematically infeasible to compute A from B.

$$A \xrightarrow{\text{trap iso}} B$$

All the functional properties of ECC are depribed:

Encryption Scheme:

Choose, Elliptic group $\mathbb{E}_p(a, b)$ and generator point, $\mathbb{G} \in \mathbb{E}_p(a, b)$ such that the smallest value of n for that $n\mathbb{G} \oplus \mathbb{C}$ is a very large prime number.

Message, \mathcal{M} is encoded in to point $F_n \in \mathbb{E}_p(a, b)$

Both sender and receiver s lects r ivate key, $n_A < n$

compute public key P_A , such that $P_A = n_A \mathbb{G}$

Cipherte. voint, $P_C = [(\mathcal{KG}), (P_M + \mathcal{K}P_B)]$

(\mathcal{K} is the random integer and P_B is the public key of receiver here). **Decryption Scheme:**

Plair expoint, $P_M \leftarrow (P_M + \mathcal{K}n_B\mathbb{G}) \leftarrow P_M + \mathcal{K}P_B$ only receiver knowing private key n_B will retrieve this point, P_M by removing $n_B\mathcal{K}\mathbb{G}$.

4. MediB Luin Protocol

In this pection we present the architectural as well as the design view of our platform. **Table 1.** describes the notations that are used in the next sections.

	Table 1: Terminology table
Notation	Description
ID	ID of the User
PWD	Password of the user
UD	Encrypted user data
U _{id}	Block id, where user data will be eved
ID _X	ID of the User X
PWD _X	Password of the user X
U _{DX}	User X's Encrypted data
U _{idX}	Block number, where user X's a. 'a is saved
Secured channel	Obtained by the authentical ons process of our system
$\mathcal{T}(\delta_{\mathrm{n}})$	Transaction of δ_n through wart contract
$\mathcal{H}_{\mathcal{M}}$	Set of all identical harmes
Γ	Address of the issuer
ν	Address of the me. sag 2 sunder
δ_{n}	Number of cat ories in the smart contract
$\{S, \mathcal{R}\}_{\texttt{authenticated}}$	authenticated set. 1e., S and receiver, \mathcal{R}
$\{S, \mathcal{R}\}_{\texttt{authenticated}}$	Unauthenti, $\mathcal{A}_{\mathcal{A}}$ tied, \mathcal{S} and \mathcal{R}
$\mathbb{B}_i, \xi_i \& \mathbb{H}_i$	Property of diffe. ant blocks

4.1. Overview of Our Protoco.

Fig. 4. shows the hig' level view of our platform. The following entities and their roles are described by there.

Data sender is the patient, who will send her personal healthcare data to the system. Data sender (a_3) s the vital role in case of data preservation. Data that will be sent to the system must be accurate otherwise wrong data will be detrimental for patient because the whole treatment depends on this sensitive data. However, our system will take the encrypted data from the users. Encryption of data will be done in the very c_{20} aning of MediBchain's process execution.

Data rece. ver will request for the data after authenticating itself to the system.

Regis tration Unit will act as an authenticator. When any party (Sender or Receiver) $x^{(1)}$ come for the first time to take the service of the system; it will store the r $y \cdot ad$ PWD to be used further. Each party will have to register for once



F[;] gure r: Hⁱgh level view of this system.

and need to preserve .nc ^D and PWD. Further they just have to log in and access through secured channel for transaction of their private data in the cloud.

Private Accessible Init (PAU) Both the parties of the system will be able to interact with FAU after authentication. It needs a secured channel to interact with PAU because dim use this unit they will send their private data to the System. It is the interm anary unit for both the levels of our system, through which the element of one level will anteract with the other.

block chain will hold the data of the users. Each transaction in the blockchain will return a_i identifier. Transaction identifiers will help the users to access the dc_{i} further.

For better understanding our system is divided into two levels. Let all is Graphical User Interface (GUI). User will interact with our system the rugh this level. Elements of level-1 are: Registration Unit and PAU. PAU is the element of both Levels so it will work between level 1 and 2. Level-2 is the broker do four system, which interacts with low level elements of this system through "AU. Element of level-2 is: blockchain. blockchain is being used as a repository of healthcare data in our system. Our platform uses permissioned blockchain which will require authentication to access.

Steps in the system : Steps of our system could be defined from Fig. 4.

Step-1 : Data sender will request with the ID and PWL to have access in the system.

Step-2: Upon accessing the system in step-2, Date sender will send data to PAU for storing.

Step-3 & 4: Step 3 & 4 will take place in level-2 of our system, where PAU will send U_{ID} to blockchain and it will return U_{IT} for future access to the blockchain and also for finding the exact Block v return data were saved.

Step-5: In this step PAU will return the O_{12} to Data sender which was given by blockchain.

Step-6 : From this step rest of the step are related to Data receiver. As step-1, this step also requires sign in r_{122} ress and after sign in Data receiver can request for the data.

Step-7: In this step Data receiver vill request for the data to Private Accessible Unit along with the U_{ID} . 'AU will receive the U_{ID} for further use.

Step-8 & 9 : Step 8 & 9 a. sar le as step 3 & 4 but the data are not same for this steps. In step-8 PAU v ill request the blockchain along with the U_{ID} and in Step-9 blockchain will return it.

Step-10 : This is ne inal step where PAU send the private data to the Data receiver.

4.2. Formal 'vesc iption of Protocol

In this section we will define how Data sender, Data receiver, and our system will work altoge her in case of sending and receiving the data. In case of data transmission in car system parties need to go through a step called registration. After confirmation of the Registration Unit that party can access the PAU.

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Figure 5: Low level view of Sending Protocol

4.2.1. Protocol between 'Jate sender and System :

Fig. 5. Shows the low x 'el /iew of sending protocol. A patient will play the role of a data sender ', this protocol. Encrypted data will be sent to the system. Generation of ciphertexts collely depend upon a function known as encryption function. General'zeo form of this function is Enc(x,y). Below we will see how this function work.

$$Enc(key, Data) = U_{\rm D} \tag{1}$$

By providing κ_{1} and the health data to this function data sender will get U_D and will send i to the system. Public key encryption technique (e.g., Elliptic Curve Cryptography (EUC)) will be applied to encrypt the private data.

Suppose X is ... Data sender of our system. At first X will request for getting into the system by providing the ID_X and PWD_X . Our system will send confirmation to X if she provides the right ID and PWD. If X could sign in to the system properly at t_{D_X} to PAU through a secured channel. Secured channel will provide the security to the transmission of data . In this stage PAU will interact with blockchain and this interaction with the blockchain will be done by the smart contracts of our system.

In our system smart contracts have been designed in a way such that blockchain will return the number of that block which has been denoted at U_{id} . Each block has a unique id which will work as an id of a specific r atient. PAU will get the U_{id} on each transaction of data in the system for X it will be $U_{ir's}$. PAU will send the U_{DX} to the blockchain then smart contract will rr such the special id U_{idX} , for X. After that PAU will send the U_{idX} to X and end the protocold. X has to store this U_{idX} otherwise next time X will not be able to acce. The rrot cond private data.

Getting the U_{idX} is the confirmation for X that means the data has been kept to the system and then X could log out and end the secured nannel transmission with the system.

4.2.2. Protocol between Data Receiver and Sys. m:

Receiving in our system will take two larces of authorization. Because after registering or signing into our system parts s will have to provide the U_{id} to get their data back through the secured changed in this phase if they fail to submit the U_{id} then they will not be able to use their data. U_{id} is the key to receive the actual data. Fig. 6. shows a low level view of receiving protocol.

Suppose user X wants to retrieve her data which she sent to the system in sending phase. As like sending phase this hase is also controlled with the authentication or Registration unit where X has to sign in first then will be able to access our system. This sign in lequires the ID and PWD of the user which was given in the registration phase. If X r rovides appropriate ID and PWD only then the system will send confirmation. After getting the confirmation X will be able to interact with the system unrough a secured channel. In this interaction with the system, X has to prove le her U_{idX}. After getting the U_{idX} system will interact with blockchain. This interaction will take place in level-2 of our system. Only PAU can interact with block chain, here the smart contracts of our system will be the medium.

Smart cor tract will send the U_{idX} to blockchain for retrieving the data of X from it. 256 bit hash *ci* the corresponding block number will be checked in the smart contract, when the hash will be matched with any block then it will continue the process to review the data. Otherwise this exception will be handled through the smart contracts.

Su voo - the hash of any block is,

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Figure 6: Low level view of Receiving Protocol

0xe3b1c14298fc1c149[,].fb⁴4c8196fb92427ae41e4649b934ca495991b785

Only if the hash of U_{TX} 's corresponding block is same then X will be able to get her data. In our system, blockchain will return the U_{DX} to PAU and it will be redirected to X lat *c*. \downarrow fter this data retrieval session will have it's end.

X will get her U_{D_A} w' ich has to be decrypted to get the actual raw data to decrypt the data user n' ed to $u_{D_A} Dec(x,y)$ function.

$$Dec(key, U_{\rm D}) = plaintext$$
 (2)

X will us equate n 2 with key and U_{DX} to retrieve the raw data.

4.2.3. storage of our system :

O r syste n will store the ID and PWD for authentication and response purpose. Our system solely will manage these private data in the cloud without deper dir 5 on any other trusted third party (TTP) apart from the PAU. Each time when user will store the data she will get a new block to write so the block-id will change by time. ID and PWD is dependent on party but U_{id} is dynamic with each data storing process.

4.3. Programmatic view of MediBchain

Smart contract of our system has been presented in this pape, through some algorithms. These algorithms have been designed to be concerted in any blockchain based language (e.g., Solidity, Golang). Contracts of our system are written in Solidity language and all the results of this paper are algorithms based on Solidity based environment. Algorithms 1-3 will be appropriate for any environment designed for blockchain environment.

Algorithm-1 describes how our system will check the *i* success verifiability. All the hash of our system is denoted by \mathcal{H}_M and all the training \mathcal{H}_i must be a part of \mathcal{H}_M . Here, *i* refers to particular number of \mathcal{H}_i .

$$\mathcal{H}_{\mathcal{M}} \leftarrow \{\mathcal{H}_1, \mathcal{H}_1, \mathcal{H}_3, \dots, \mathcal{H}_i\}$$

 $\mathcal{H}_{\mathcal{M}}$ is the set of all identical hashes of our system that will be provided in the time of account creation in the block of a network. $\Gamma \& v$ are part of $\mathcal{H}_{\mathcal{M}}$ and play significant role in transaction. Two different notations have been used to reduce the complexity of Algorithm 1, issuer of contract has been denoted with Γ and data uploader/downloade has been denoted with v. Here, issuer is the address who runs the contract and message lender is the address who sends the message. If both of them are not so as t in Algorithm 1 will return false.

Algorithm 1: Checking on Tssher and Sender		
Result: Verified 7, 55. °r		
1 Γ, ν;		
2 ▶ address of the issuer and message sender respectively		
3 while $\underline{\Gamma} \&\& v \in \mathcal{A}_M$ do		
4 if $\Gamma \neq th' n$		
5 reu."r,		
6 els		
7;		
8 ► It will proceed the code to next algorithm		
9 end		
10 en."		

In-, ... gorithm is important for security and accountability of data transaction.

It will work in between the time of smart contract execution and be data transaction (e.g, upload, download) between MediBchain and blocker. in. L. vesdroppers could take a chance of data manipulation in the meantim. All up accounts of this system will be the part of \mathcal{H}_M and also the initiator of contract and data uploader/downloader will be same. Execution of rest of the contract will be dependent on the similarity of $\Gamma \& \nu$. Algorithm-2 will be initiated after algorithm-1, in which δ_n represents the number of categories to be held by the structure of data in our contracts.

Alg	Algorithm 2: Transaction of Data	
R	Result: Data Upload	
1 S	1 struct Data $\leftarrow \sum_{1}^{n} \delta_{n}$	
2 D	2 Data[] data;	
3 b	ool←0;	
4 while <u>n</u> do		
5	▶ getting input from message sender, v	
6	if <u>v</u> returns string then	
7	$data \leftarrow \sum_{1}^{n} \delta_{n};$	
8	bool $\leftarrow 1$;	
9	return bool;	
10	else	
11	return bool;	
12	12 end	
13 e	nd	

Algorithm-2 will be *crec* ated after fulfilling the conditions below. Iff,

 Γ && $\nu \in \mathcal{H}_{\mathcal{M}}$ and, Γ (issuer) = ν (message sender)

Here, Γ is the *i* ddre. who runs the contract and *v* is the address who sends the message. If b*i* th *c*, them are not same then this algorithm will return false. Users' (patients') data fill *i* e having different categories to be inputted. Different categories mean that the healthcare data come in different types, suppose user wants to save B bod subars data and also Blood pressures data these two are different. By category the refer to this scenario that the user can store different diagnostic results in a block. Hence, we have designed two different contracts. In Algorithm-2 each structure will hold maximum four different types of healthcare data to be stored in the block if we change data part as follows,

 $data \leftarrow \sum_{n=1}^{4} \delta_n$ We have another smart contract which takes maximum eight ω even, types of health data to be stored in the block. For that we need to charge data part again. So above data part will be changed as follows-

$$data \leftarrow \sum_{n=1}^{8} \delta_n$$

We have shown some computational analysis in subsect on- 3.3 using the variation of data storing capabilities of different smart contrac.s.

Line-1 is showing that the structure of smart contrac' can take n number of individual data from a particular patient at a time. In the $\log \rho$ da a will be assigned to its corresponding structure in line-7 and then the a. 'a will be written in the block in the same contract. A particular structure will be writen in a particular block. As mentioned earlier each block of blockchain hour different id which is not same as $\mathcal{H}_{\mathcal{M}}$. $\mathcal{H}_{\mathcal{M}}$ represents the account id of block digin network whereas hash ids has been denoted with \mathbb{H}_i . A bool variable was been returned from Algorithm-2 as a flag for Algorithm-3. In Algorithm-3. ξ_i rep. sents the block number and \mathbb{H}_i represents the hash of particular block.

Algorithm 2. Diock id Consection		
Algorithm 3: Block-id Generation		
Result: Block-id		
1 ξ_i , \mathbb{H}_i ;		
2		
3 while bool do		
4 ▷ Returned value from Algorithm 2		
5 if $\underline{bool} \leftarrow 1$ ther		
$6 \qquad \xi_i \leftarrow block Nu. ber);$		
7 $\mathbb{H}_i = \text{bloc}^1$. '¬lockhash(ξ_i);		
8 return \mathbb{H}_i		
9 else		
10 return $1^{\tau_1\tau}$;		
11 end		
12 end		

Algori am-3 will return hash id \mathbb{H}_i if all the requirements will be fulfilled by the co. tract. : will take a variable named *bool* by which this algorithm will define 'hetner' to return block-id, H_i or not. Functions block.Number() and bloc :.bloc':hash() are the syntex of Solidity language, where block.Number() will re. $rn t^{\dagger}$ e corresponding block number ξ_i and block.blockhash() will retu 🗤 📰



 ξ_i and \mathbb{H}_i are the properties of each block, \mathbb{B}_i by which our system will work

 $\mathbb{H}_i \leftarrow block.blockHash(\xi_i)$

Programmatically each \mathbb{H}_i will be generated from it's corresponding ξ_i . As instance, if block.blockhash() gets ξ_1 as a parameter it vill return \mathbb{H}_1 or if it gets ξ_{20} the function will return \mathbb{H}_{20} . So the relation can be written as,

 $\{\mathbb{H}_1, \mathbb{H}_2, \mathbb{H}_3, \mathbb{H}_4, \dots, \mathbb{H}_i\} \equiv \{\xi_1, \xi_2, \xi_3, \xi_1, \dots, \xi_i\}$

5. Protocol Analysis & Evaluation

5.1. Security Analysis

- **Pseudonymity**: Data Sender, S and Feceiver, R will not be identified by any party during transaction.
 - Pseudonymity of S: A for an entication S will upload the encrypted private data, U_D . Any our party will not be able to identify S by looking her U_D because of it's identificationless attribute.
 - Pseudonymity of R: H_i will be used to trace particular B_i of the blockchain which holds the private data of S. During transaction T party will hold the H_i to have her U_D back from the system, these H_is are as sensitive with private data for receiver. H_i will be held by only our party which ensures the pseudonymity of Data Receiver because no one will be while to detect S during T or even after T because of encrypted property of U_D.

Support x{ID,PWD} is the function for authentication,

$$\alpha$$
{ID,PWD} \longrightarrow { S, \mathcal{R} }_{authenticated}

- **Privacy:** R gistration Unit and U_D ensures the privacy of the $\{S, \mathcal{R}\}_{authenticated}$ and data respectively.
 - P.ivacy from system: Parties, $\{S, \mathcal{R}\}_{authenticated}$ of our system have privacy as pseudonymity of users is maintained. α {ID,PWD} will ensure

the access in the system. This controlled access of $\{\mathcal{L},\mathcal{R}\}_{uthenticated}$ provide privacy to the users of our system. Therefore, $\mathbb{C}_{\mathcal{T}}$ of $\{\mathcal{S},\mathcal{R}\}_{authenticated}$ can not be compromised any way.

• Privacy from other parties: S will have her dedica. $d^{T} J_{i}$ in the blockchain to store U_D. So, if any $\{S, \mathcal{R}\}_{authenticated}$ of or _ysten, tries to access any other party's data it will not be able to access the particular block as each party will have their dedicated \mathbb{H}_{i} .

Clearly, the former analysis guarantees a very strong privacy of parties because only $\{S, \mathcal{R}\}_{authenticated}$ will be able to crees. As well as retrieve data from that particular \mathbb{B}_i .

• Integrity:

- Access request data integrity Each time S or R tries to access the system, she needs to authenticate burself primarily. This access request needs to be done by both and dy amic entities- S and R of system. These access requests will require correct ID and PWD, which will be generated by party itsen and will be holding by the database of system. So without S or R and system these authentication data will not be known by anyone. Ly which system guarantees the access request data integrity.
- User data Integrity Use of *Enc(x,y)* function ensures the data integrity as the data in *the* blockchain will make no sense to any other person except the data owner. After retrieving the data from the system {S,R}_{authentica} a need to decrypt the U_D with *Dec(x,y)* function. In order transfer with integrity level attacker needs to break the security of un orlying encryption scheme, ECC.

All the i ata that are related to our healthcare data management system guarantees in 2 (rity

• Accumtalility:

• T ansactional B_{*i*}: When any party will come to save it's data to the system a unique number or nonce, H_{*i*} will be returned which leverages the accountability of our system. Only party itself will be holding this

nonce which makes the party accountable for it's U_D $\ensuremath{\mathsf{L}}\xspace$ without valid information about

 $\mathbb{B}_i \ni \xi_i, \mathbb{H}_i$ party will not be able to access her _F ivate data from blockchain.

- PAU as bridge: Interaction of $\{S, \mathcal{R}\}_{authentic \ ced}$ which the system is controlled. This controlled path refers to the secure 1 channel which will be created by the party itself through $\alpha\{1, \mathbb{P}W'D\}$. Through this channel $\{S, \mathcal{R}\}_{authenticated}$ will interact with $\mathbb{P}^{A'}D'$ which is a bridge between the system and blockchain. Secure ' char hel makes the bridge accountable for secured \mathcal{T} with blockchain.
- Security: Each \mathbb{B}_i will be dedicated to $\{\mathcal{C}, \mathcal{R}\}_{\text{authenticated}}$ and their \mathbb{H}_i is secured as integrity is guaranteed in \mathbb{C} ." platform. As a result, these \mathbb{B}_i will not be accessed by any $\{S, \mathcal{R}\}_{\text{auth}}$. If attacker somehow manages to intrude into the blockchain netvor'. patients' sensitive data will make no sense because of encrypted "tribu." of data. Accessing the raw data of patient will need the keys and $D_{\mathbb{C}}(x,y)$ will return the raw data to parties. So, the data security is guaranteed.

The equation for Transactior

 $\mathcal{T}(\delta_{n}) \longleftarrow \left\{ \left\{ \forall \mathcal{H}_{\mathcal{M}} : \Gamma \in \mathcal{H}_{\mathcal{M}}, \forall \in \mathcal{H}_{\mathcal{M}}, \Gamma = \nu \right\} \text{ and } \left\{ \forall \alpha_{\{\mathcal{S},\mathcal{R}\}_{\text{authenticated}}} \{ID, PWD\} \right\} \right\}$

After analyzing each of the p operties we can conclude with saying that no platform secures blockchain build preudonymous healthcare data other than our platform-'MediBchain', in the 1 est of our knowledge.

5.2. Computation - Tvaluation

We setup an Civir onment to evaluate our protocol by writing programs using Solidity 0.4.11 and JATA 1.8 with a computer Intel(R) Core(TM) i5, CPU-3.30 GHz, 8 GB of RAM, Win 8, 64-bit OS. We deployed Elliptic Curve Cryptography (ECC) for generatin, and retrieving the input and output respectively.

5.3. Dana sharing:

W test t e computation time to generate the cipher texts. Each encryption is an isolated process. Fig. 7. presents the data encryption time versus string size of nealthcare data. We take several inputs to see how the rate of growth of



Figure 7: Computation in generating input

time for encryption changes with variable input size. We take 5 to 30 kilobytes of data to analyze the encryption time of different data size. From the resultant graph we can observe the rate of grow b of curve is nearly linear which means the encryption time increases with increase of data size. Data sharing phase of our system is variable and independent process, variable means that input size could vary for different users and independent means the encryption of different users' data are not dependent on each other.

5.3.1. Data manipulat on will mart contract:

Before $stin_{\mathcal{S}}$ access of a block in the blockchain network data will be accessed by our smart contract. Use of smart contract will cost some gas which is known as the cry, toful of Ethereum Virtual Machine (EVM). To run any dapp (distributed apprection) on the Ethereum environment the executed application will need to



Figure 8: Computation cost of transaction and vecution of smart contract

have some transactions in the networ, in return of transaction the environment costs the executor some gas. Initiator of executor of transaction will get the gas in exchange of Ether in Ethereum environment. We evaluate two smart contracts, one with 4 inputs category other with 8 inputs category. In context of programming language which is number of variables to take input from party. Subsections 5.3.2 and 5.3.3 will depict the analogy of different terms of smart contract with 4 inputs category and 5.3.4 and 0.3.5 will depict the analogy between two different smart contracts with the transaction and the inputs, where the transaction and the inputs category based on the transaction and the cost of our smart contract.

5.3.2. Transaction Cost vs Execution Cost:

Fig. 8. depict. t^{1} e analogy between transaction and execution cost of smart contract. To t^{1} are an accurate analyzing result we run the smart contract with different input sizes t^{1} at varies from 5 to 100 characters of string. Curves in Fig. 8. shows the cost t_{1} increasing with the input size. But the rate of growth of these two curves is san t^{2} between the intervals and linear too.

5.3.3. Block-id generation costs:

O e of the key terms to be ensured while writing smart contracts was block-id generation. Block-id generation will cost for execution and transaction. We analy e the clock-id generation cost with different string length, but interestingly it



Figure 9: Computation cost of transaction and Cocution of smart contract

costs same for all the inputs. Fig. 9. shows the curves of execution and transaction cost of block-id generation. It is clean that each parameter is almost constant with the increase of the size of string. Transaction and execution cost is same for growing input size.

5.3.4. Transaction cost of variable inpu.::

Parties of our system may i_{10} to upload a vast amount of data in different categories. Smart contract may have to be redesigned so that we analyze the cost to see how our platform reacts with an increasing amount of category to store it in blockchain. Before this subsection we were talking only about smart contract having 4 categories to fake with put, but for having an effective analogy we will give 8 categories as in put to see how the behavior changes of our platform. Fig. 10. shows us the analogy we tween two smart contracts in which one will take 4 inputs and other will ake 8 inputs. In Fig. 10. we can see that smart contract having 8 categories of input will cost higher, but the rate of growth of curves are similar and the cost with increase with string size.

5.3.5. Execution . 1 of variable inputs:

Fig. 1., presents the execution cost of smart contract with variable input. As explained those the mart contracts may vary in different scenario, so that we present the extraction costs' analogy in Fig. 11. The rate of growth of curves is similar but smart contract with 8 inputs will cost more gas while execution with increasing string lenguls.



Figure 10: Transaction cost of smart contract with variable input

5.4. Output generation:

To get the plaintext or actual private her theorem data of patient the data from blockchain need to be decrypted. As his enclyption, decryption or output generation process is also isolated. All the output generation for the parties is independent from each other. To analyze the output retrieval time we take different sets of string 5 to 30 kilobytes of data at a single input to get an actual idea of output retrieval time for the patients. In high 12, curve shows that the rate of growth of time is related with the input size as the time is increasing for decryption with input size. The curve is rearly linear. Time is in millisecond in the graph that is computed with Java dering de ryption. Elliptic Curve Cryptography (ECC) is used to generate the plaintext.

5.4.1. Input generation vs catput retrieval:

Generation of inp t and output is independent from each other. Encryption will take place in u. time of giving input and decryption will take place in the time of output Fi₂. 13. depicts that two processes take very different amount of time while proc_2 sins. With the string length both the time increase but encryption needs more than decryption. For encryption it takes 80 to 90 milliseconds where decryption needs less than 10 milliseconds.

6. C(nclusi, n

The p_{mp} er presented privacy preserving platform for healthcare data in cloud. We have defined a set of security and privacy requirements for healthcare data



Figure 11: Execution cost of smark ontract with variable input

management systems and argued why such attributes are necessary for a healthcare data management system in cloud. Our analysis shows that our platform satisfies all such requirements. Experimental performance evaluation shows that this platform runs well in blockchain environment. In the future we will try to explore the interoperability between different entities (e.g., diagnostic center, hospital, doctors, patients) of hering reprocess, and another direction would be to address the issue of handling key-th ft/loss mechanisms or key distribution techniques.

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Figure 12: Computation time in genearing actual output

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Figure 13: Input generation vs Jun Contrieval time of system

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Highlights:

- 1. We proposed a user centric EHR systems for Healthcare, which gives the total controlling power of the data to the users.
- 2. In generic EHR platform, it becomes easier target of intruders to new ide the system than totally breaking the security of the system. We solved this problem by implementing permissioned Blockchain along with the cryptographic function.
- 3. We explored the archival use of Blockchain in our platform by storing the data of users in the blocks of the permissioned Blockchain.
- 4. Controlling the pseudonymity of the users is a big challenge. V a solved the pseudonymity issue by applying cryptographic function. We used Elliptic Curve Cryptography (ECC) to make the data safe from other party in this distributed ledge system.